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Introduction

The *Threading Building Blocks* (TBB) is a C++ library developed by Intel to specifically address programming of multi- and many-core systems. It supports Linux, Windows and MacOSX and all major C++ compilers.

TBB is available as a commercially supported library from Intel or as an open-source version from

http://threadingbuildingblocks.org

TBB provides algorithms and data structures to define tasks in a parallel program. These tasks are then mapped by an internal scheduler to worker threads.

In contrast to other thread programming packages the programmer has no access to these threads, only to the tasks.

TBB is using C++ templates extensively to minimise runtime overhead.

Remark

*This course is based on Version 4 of TBB.*
Scheduling

By default tasks are enqueued into thread-local work queues.

In case of a load imbalance, *task-stealing* is used to transfer tasks to other threads.

Also, the scheduler tries to preserve *cache locality* by scheduling tasks first, which have been most recently in the cache.

This may result in *unfair* scheduling, i.e. other tasks may starve for processing resources.

Limitations

Due to unfair scheduling, TBB is *not* designed for

- I/O operations, where a task may block until data can be fetched or
- realtime operations, since no guarantee upon the execution time can be given.
The standard introductory example looks as follows using TBB:

```cpp
#include <iostream>
#include <tbb/tbb.h> // include all of TBB
using namespace tbb; // import TBB namespace

struct hello : public task { // Hello World Task
    task * execute () {
        std::cout << "Hello, world!" << std::endl;
        return nullptr;
    }
};

int main () {
    hello & t = * new( task::allocate_root() ) hello;
    task::spawn_root_and_wait(t);
    return 0;
}
```

Here, printing “Hello, World” is defined as a task and handed to the TBB scheduler for execution.

Using the Intel compiler, a command line switch may be used to enable TBB:

```
> icpc -tbb -o hello hello.cc
```

All other compilers have to use the standard flags for include and library paths:

```
> g++ -I<path to TBB>/include -L<path to TBB>/lib -ltbb -o hello hello.cc
```
The TBB library needs no explicit initialisation. By default it will create worker threads for each processor (core).

To initialise the TBB library with an explicitly defined number of worker threads, the class `task_scheduler_init` is provided:

```cpp
#include <tbb/task_scheduler_init.h>

void
main () {
  tbb::task_scheduler_init tsi( nthreads );

  ...
}
```

**Remark**

*The worker threads include the main thread, e.g. the initial thread of the process.*
Lambda Functions
Lambda Functions

Lambda functions are *anonymous* functions in C++11, also called *closures*, which are especially useful when using TBB.

A lambda functions consists of a function body and a data environment, similar to a task.

The (almost) full definition of lambda functions is

\[
\left[ \text{capture} \right] \left( \text{params} \right) \rightarrow \text{ret-type} \{ \text{body} \}
\]

With

- **capture**: defining the data environment,
- **params**: function parameters,
- **ret-type**: defines the function return type and
- **body**: is the function body.

The last three correspond to the definition in standard named functions.
Captured Variables

Variables referenced in the function body have to be imported to the data environment of the lambda function.

Variables can either be captured by *copy* or by *reference*.

To capture a variable by copy, the variable has to be explicitly listed in the capture description:

```
[ x, y, i, j ] ...  // capture copy of x, y, i, j
```

Alternatively, with

```
[ = ] ...  // all captured by copy
```

*all* variables referenced in the function body are automatically captured by copy.

If a variable should be captured by reference, the variable has to be prefixed by *&* in the capture description:

```
[ & x, & y, & i, & j ] ...  // capture reference of x, y, i, j
```

To capture all variables in the function body by reference, use

```
[ & ] ...  // all captured by reference
```
Captured Variables

Capture by copy and by reference may also be combined:

```
[ x, & y, i, & j ] ... // capture x, i by copy and y, j by reference
```

or

```
[ =, & y, & j ] ... // capture all by copy except y and j
```

Remark

When using “=”, it has to be first in the capture description.

An empty capture description will not capture any variable:

```
[ ] ... // do not capture any variable
```
Captured Variables

The exact value of a captured variable when executing the lambda function depends on how it is captured:

```cpp
int i = 1, j = 2;
double x = 3, y = 4;

auto f1 = [x, y, i, j] () -> void { std::cout << i << j << x << y << std::endl; };
auto f2 = [=] () -> void { std::cout << i << j << x << y << std::endl; };
auto f3 = [&x, &y, &i, &j] () -> void { std::cout << i << j << x << y << std::endl; };
auto f4 = [&] () -> void { std::cout << i << j << x << y << std::endl; };
auto f5 = [x, &y, i, &j] () -> void { std::cout << i << j << x << y << std::endl; };
auto f6 = [=, &y, &j] () -> void { std::cout << i << j << x << y << std::endl; };

i = 5; y = 6;
f1(); f2(); f3(); f4(); f5(); f6();
```

The output is then:

```
1234
1234
5236
5236
1236
1236
```

Remark

Lambda functions have no explicitly nameable type. But it can be inferred by `auto`.
The full lambda definition can be simplified to

```
[ capture ] ( params ) { body }
```

if the body contains a single return statement, from which the return type can be automatically inferred:

```cpp
auto f = [] ( const double x ) { return x * x; };
```

If no return statement exists, the return type defaults to `void`:

```cpp
auto f = [] ( const double x ) { compute_something( x ); };
```

Similarly, if the lambda function has no parameters, it can be simplified to

```
[ capture ] -> ret { body }
```

e.g.

```cpp
auto pi = [] -> double { return 355.0 / 113.0; };
```

Both simplifications can also be combined to

```
[ capture ] { body }
```
Loops
The parallelisation of simple loops, e.g.

```
for ( size_t i = 0; i < n; ++i ) {
    ...
}
```

is provided by the TBB algorithm

```cpp
template<typename index_t, typename func_t>
func_t parallel_for ( index_t start,
                     index_t end,
                     const func_t & f );
```

Using `parallel_for`, the above loop translates into

```
parallel_for( 0, n, [] ( index_t i ) { ... } );
```

Alternatively, the loop body can be encapsulated by a named function:

```cpp
void f ( index_t i ) {
    ...
}

parallel_for( 0, n, f );
```
**Example: Matrix Multiplication**

Computing $C = A \cdot B$ with $A, B, C \in \mathbb{R}^{n \times n}$ in parallel using `parallel_for`:

```cpp
parallel_for( 0, n,                      // parallelise outer most loop
    [n,&A,&B,&C] ( const size_t i ) {  // capture A,B,C by reference
    for ( size_t j = 0; j < n; ++j ){
        double f = 0;
        for ( size_t k = 0; k < n; ++k )
            f += A(i,k) * B(k,j);
        C(i,j) = f;
    }
});
```

The parallel speedup of the above implementation for $n = 1024$ is:

<table>
<thead>
<tr>
<th></th>
<th>Xeon X5650 (24 threads)</th>
<th>Xeon E5-2640 (24 threads)</th>
<th>XeonPhi 5110P (240 threads)</th>
</tr>
</thead>
<tbody>
<tr>
<td>outer most loop</td>
<td>11.26x</td>
<td>11.29x</td>
<td>186.05x</td>
</tr>
</tbody>
</table>
The parallel_for algorithm also accepts a `step` parameter:

```cpp
template<typename index_t, typename func_t>
func_t parallel_for ( index_t start,
                      index_t end,
                      index_t step,
                      const func_t & f );
```

which enables the parallelisation of loops of the form

```cpp
for ( size_t i = 0; i < n; i += step ) {
    ...
}
```
The `parallel_for` algorithm also accepts a `step` parameter:

```cpp
template<typename index_t, typename func_t>
func_t parallel_for ( index_t start,
                      index_t end,
                      index_t step,
                      const func_t & f );
```

which enables the parallelisation of loops of the form

```cpp
for ( size_t i = 0; i < n; i += step ) {
    ...
}
```

**Remark**

*The loop must not wrap around and the step value must be positive.*
Simple Loops

Using the above form of `parallel_for`, each iteration of the loop induces a new task, i.e. a fine granular approach to loop parallelisation.

If the work per iteration is small, the management overhead may be too big for efficient parallelisation.

For such cases, the generalised version of `parallel_for` can be used, which uses `ranges` to specify the index set to operate on:

```cpp
template< typename range_t, typename body_t >
func_t parallel_for ( const range_t & range,
                      const body_t & body_t & body );
```

The `body` object may be a lambda function

```cpp
[...]( const range_t & r ) { ... }
```

or a standard C++ object with the following minimal interface:

```cpp
struct body_t {
    body_t ( const body & );        // copy constructor
    ~body_t ();                     // destructor
    void operator () ( const range_t & r ) const; // range based operator
};
```
TBB uses *ranges* to partition a given index set into sub sets for task definition. For this, it recursively subdivides the index set, until a indivisible set is reached. The constructed tasks are assigned to (or stolen from) the worker threads.

Each range in TBB has to implement a minimal interface to enable partitioning:

```cpp
struct range_t {
    range_t ( const range_t & ) // copy constructor
    range_t ( range_t & r, split ) // split range into two sub ranges
    ~range_t (); // destructor

    bool empty () const; // return true if range is empty
    bool is_divisible () const; // return true if range can be partitioned into
                                 // two sub ranges
};
```

The *splitting* constructor partitions the given range into two sub ranges, usually of equal size, and updates the range parameter with one of the new sub ranges:

```cpp
struct int_range_t {
    int lb, ub; // lower and upper bound of interval [lb,ub)

    int_range_t::int_range_t ( int_range_t & r, split ) {
        int mid = (r.lb + r.ub) / 2;
        lb = mid; ub = r.ub; // this becomes second sub interval
        r.ub = mid; // r becomes first sub interval
    }
};
```
TBB defines ranges for one, two and three dimensional indexsets.

```
template < typename value_t > blocked_range;

template < typename row_value_t,
type name col_value_t> blocked_range2d;

template < typename page_value_t,
type name row_value_t,
type name col_value_t> blocked_range3d;
```

For `value_t`, any integral or pointer datatype may be used. Furthermore, standard STL `random` iterators, e.g. for `std::vector`, are possible.

```
blocked_range< size_t > r1( 0, 100 ); // index set [0,100)
branked_range2d< size_t > r2( 0, 100, 0, 10 ); // index set [0,100) x [0,10)
branked_range3d< size_t > r3( 0, 100, 0, 20, 0, 5 ); // index set [0,100) x [0,20) x [0,5)
```

Range iteration is performed in the standard STL way using `begin()` and `end()`. For the two and three dimensional ranges, the corresponding sub ranges are accessed using `rows()` and `cols()` and, for `blocked3d_range` via `pages()`:

```
for ( auto i = r1.begin(); i != r1.end(); ++i ) ...

for ( auto i = r2.rows().begin(); i != r2.rows().end(); ++i )
  for ( auto j = r2.cols().begin(); j != r2.cols().end(); ++j ) ...

for ( auto k = r3.pages().begin(); k != r3.pages().end(); ++k )
  for ( auto i = r3.rows().begin(); i != r3.rows().end(); ++i )
    for ( auto j = r3.cols().begin(); j != r3.cols().end(); ++j ) ...
```
Using ranges, the parallel loop

```cpp
parallel_for( 0, n, [] ( index_t i ) { ... } );
```

can be rewritten as

```cpp
parallel_for( blocked_range< size_t >( 0, n ),
             [] ( const blocked_range< size_t > & r ) {
                for ( auto i = r.begin(); i != r.end(); ++i ) {
                    ...
                }
            } );
```
Using ranges, the parallel loop

```cpp
parallel_for( 0, n, [] ( index_t i ) { ... } );
```

can be rewritten as

```cpp
parallel_for( blocked_range< size_t >( 0, n ),
             [] ( const blocked_range< size_t > & r ) {
               for ( auto i = r.begin(); i != r.end(); ++i ) {
                 ...
               }
             } );
```

**Example: Matrix Multiplication**

With `blocked_range2d`, both outer loops can be parallelised:

```cpp
parallel_for( blocked_range2d< size_t>( 0, n, 0, n ),
              [n,&A,&B,&C] ( const blocked_range2d< size_t > & r ) {
                for ( auto i = r.rows().begin(); i < r.rows().end(); ++i ) {
                  for ( auto j = r.cols().begin(); j < r.cols().end(); ++j ) {
                    double f = 0;
                    for ( size_t k = 0; k < n; ++k )
                      f += A(i,k)*B(k,j);
                    C(i,j) = f;
                  }
              } } );
```

This yields a speedup of:

<table>
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<th>XeonPhi 5110P (240 threads)</th>
</tr>
</thead>
<tbody>
<tr>
<td>both outer loops:</td>
<td>11.20x</td>
<td>11.29x</td>
<td>211.88x</td>
</tr>
</tbody>
</table>
Range Concept

Beside the lower and upper bound for each dimension, an optional *grain size* may be specified, which defines the minimal chunk size for recursive partitioning, which is *one* by default.

```c
blocked_range ( value_t lb, value_t ub, size_t grainsize );
blocked_range2d ( value_t row_lb, value_t row_ub, size_t row_grainsize,
                 value_t col_lb, value_t col_ub, size_t col_grainsize );
blocked_range3d ( value_t page_lb, value_t page_ub, size_t page_grainsize,
                 value_t row_lb, value_t row_ub, size_t row_grainsize,
                 value_t col_lb, value_t col_ub, size_t col_grainsize );
```

In the following example, the task borders for different grain sizes for a two dimensional index set are shown. Hereby, the uneven one dimensional range sizes are a result of the recursive sub division.

```c
blocked_range2d< size_t > r1( 0, 10, 2
  0, 20, 3 );
blocked_range2d< size_t > r2( 0, 10, 4
  0, 20, 6 );
```
For the computation $x_i = x_i + \sqrt{y_i}$ with $x, y \in \mathbb{R}^n$, the parallel speedup depending on the grain size on a 2-CPU Xeon E5-2640 is presented in the following diagram:

![Graph](image)

For a very small grain size, management overhead clearly dominates the runtime, preventing any speedup.
Ranges define a partition of an index set, e.g. into chunks of at most grain size. However, a parallel loop starts with the whole index set and successively divides it into smaller sub index sets, such that each thread will have some tasks to work on.

This process is controlled by *partitions* in TBB, of which the following are provided:

- **auto_partitioner**: divide range into sub ranges until all threads have (almost) equal load,
- **affinity_partitioner**: similar to auto_partitioner but tries to maximise cache locality,
- **simple_partitioner**: divides range until grain size is reached.

The partitioner is an optional argument to `parallel_for`:

```cpp
parallel_for ( start, end, func, partitioner );
parallel_for ( range, body, partitioner );
```

e.g.

```cpp
parallel_for( blocked_range< size_t >( 0, n ),
    [] ( const blocked_range< size_t > & r ) { ... },
    simple_partitioner() )
```

---

Kriemann, »Introduction to Parallel Programming«
By default, `auto_partitioner` is used. Since the sub division is stopped, when all threads have enough work, the specified grain size is not necessarily reached. Therefore, even if the grain size is one, the efficiency will often *not* deteriorate.

In contrast to this, the `simple_partitioner` will always split ranges until the the sub ranges are not larger than the grain size, hence, the latter needs to be chosen *appropriately*.

The `affinity_partitioner` may increase performance if several iterations are performed for the *same* data set and the per task data *fits* in the local cache.
Parallel reduction operations are provided in TBB by the `parallel_reduce` algorithm:

```cpp
template< typename range_t, typename value_t,
          typename func_t, typename reduce_t >
value_t parallel_reduce ( const range_t & range,
                         const value_t & identity,
                         const func_t & func,
                         const reduce_t & reduce );
```

Here, `func` is the function accumulating values over a given range. It has to implement

```cpp
struct func_t {
    value_t operator () ( const range_t & r, const value_t & val );
};
```

where `val` specifies a start value for the local reduction.

The combination of locally computed values is implemented by `reduce_t`:

```cpp
struct reduce_t {
    value_t operator () ( const value_t & x1, const value_t & x2 );
};
```

Finally, `identity` denotes the identity with respect to the reduction operation.
For a standard reduction based on addition, the corresponding `func_t` implementation could be

```cpp
struct arraysum_t {
    std::vector<double> & x; // array of coefficients to reduce

    arraysum_t ( std::vector<double> & ax ) : x( ax ) {} // constructor

    double operator () ( const blocked_range<size_t> & r, const double val ) {
        double s = val;
        for ( auto i = r.begin(); i != r.end(); ++i )
            s += x[i];
        return s;
    }
};
```

Similarly, for `reduce_t`, one could use

```cpp
struct plus_t {
    double operator () ( const double x1, const double x2 ) const {
        return x1+x2;
    }
};
```

Finally, the parallel reduction looks as

```cpp
sum = parallel_reduce( blocked_range<size_t>( 0, n, OPT_GRAIN_SIZE ),
                      double(0),
                      arraysum_t( x ),
                      plus_t() );
```
The main advantage of TBB reductions over other implementations, e.g. OpenMP, is that it is *not* restricted to elementary data types.

As an example, if the scalar values in the previous summation are replaced by vectors in $\mathbb{R}^3$, the parallel reduction based on lambda functions is:

```cpp
std::vector< Vector3 > vectors( n );
...
sum = parallel_reduce( blocked_range< size_t >( 0, n, OPT_GRAIN_SIZE ),

    Vector3(0,0,0),

    [&vectors]( const blocked_range< size_t > & r,
                 const Vector3 & val ) {
        Vector3 v = val;
        for ( auto i = r.begin(); i != r.end(); ++i )
            v += vectors[i];
        return v;
    },
    std::plus< Vector3 >()
);
```

All that is needed here are operators $\text{+=}$ and $\text{+}$ for objects of type $\text{Vector3}$.

**Remark**

_The STL provides a wide variety of binary operations in the _functional_ module, e.g. plus, minus, multiplies, divides, etc._.
Extending the previous example to matrices would result in many copy operations, since each local reduction or combination would return a matrix object.

To avoid such copies, TBB provides an imperative form of parallel reductions:

```cpp
template< typename range_t, typename body_t >
value_t parallel_reduce ( const range_t & range,
                        const body_t & body );
```

Here, `body_t` must implement the following interface:

```cpp
struct body_t {
    body_t ( body_t & b, split ); // splitting constructor
    ~body_t (); // destructor
    void operator () ( const range_t & r ); // sub range accumulation
    void join ( const body_t & b ); // combine parameter with local result
};
```

The splitting constructor is used by TBB to make copies of the `body_t` object for different worker threads.

It may be executed concurrently with the sub range accumulation or the `join` function. Hence, special protection for object local variables may be necessary, although this is usually not needed.
Using the imperative form of `parallel_reduce`, the matrix sum $\sum_i M_i$ with $M_i \in \mathbb{R}^{m \times m}$, $0 \leq i < n$, is computed via:

```cpp
struct MatrixSum {
    std::vector< Matrix > & matrices;
    Matrix loc_sum;

    MatrixSum ( MatrixSum & ms, split )
    : matrices( ms.matrices ), loc_sum( ZERO_MATRIX )
    {}

    void operator () ( const blocked_range< size_t > & r ) {
        Matrix M( ZERO_MATRIX );
        for ( auto i = r.begin(); i != r.end(); ++i )
            M += matrices[i];
        loc_sum += M;
    }

    void join ( MatrixSum & M ) {
        loc_sum += M.loc_sum;
    }
};
```

Here, no additional copy is performed, since only local variables are updated and *not* returned from the involved functions.

The actual parallel reduction then looks as

```cpp
std::vector< Matrix > matrices( n );
...
MatrixSum sum( matrices );
parallel_reduce( blocked_range< size_t >( 0, n, OPT_GRAIN_SIZE, sum );
```
Example: Matrix Multiplication

Parallelisation of $C = A \cdot B$ may be extended to the innermost loop, where a reduction is performed for the dot product:

```cpp
parallel_for(
    blocked_range2d< size_t >( 0, n, OPT_GRAIN_SIZE,
                               0, n, OPT_GRAIN_SIZE ),
    [n,&A,&B,&C] ( const blocked_range2d< size_t > & r ) {
        for ( auto i = r.rows().begin(); i < r.rows().end(); ++i ) {
            for ( auto j = r.cols().begin(); j < r.cols().end(); ++j ) {
                C(i,j) = parallel_reduce( blocked_range< size_t >( 0, n, OPT_GRAIN_SIZE ),
                                          double(0),
                                          [i,j,&A,&B,&C]( const blocked_range< size_t > & ri,
                                                        double val ) {
                                              double f = val;
                                              for ( auto k = ri.begin(); k != ri.end(); ++k )
                                                  f += A(i,k)*B(k,j);
                                              return f;
                                          },
                                          std::plus< double >() );
            }
        }
    } } } );
```

Comparing the previous approaches for $n = 1024$ and a grain size of 64 we have:

<table>
<thead>
<tr>
<th></th>
<th>Xeon X5650 (24 threads)</th>
<th>Xeon E5-2640 (24 threads)</th>
<th>XeonPhi 5110P (240 threads)</th>
</tr>
</thead>
<tbody>
<tr>
<td>outer most loop only:</td>
<td>10.41x</td>
<td>9.80x</td>
<td>16.16x</td>
</tr>
<tr>
<td>both outer loops:</td>
<td>11.31x</td>
<td>11.00x</td>
<td>121.47x</td>
</tr>
<tr>
<td>outer loops plus reduction:</td>
<td>10.08x</td>
<td>10.25x</td>
<td>167.39x</td>
</tr>
</tbody>
</table>
By default, the splits and joins of parallel reduction operations may be performed in any order. Since for the reduction itself, associativity is assumed, this poses no problem \textit{in theory}.

In practise however, the results of parallel reductions may differ between different runs.

The enforce the same order of splits and joins, TBB provides a deterministic variant of \texttt{parallel\_reduce}:

\begin{verbatim}
template< typename range_t, typename value_t,  
          typename func_t,  typename reduce_t > 
value_t parallel_deterministic_reduce ( const range_t & range,  
                                      const value_t & identity,  
                                      const func_t & func,  
                                      const reduce_t & reduce );
\end{verbatim}

\begin{verbatim}
template< typename range_t, typename value_t,  
          typename body_t > 
value_t parallel_deterministic_reduce ( const range_t & range,  
                                      const body_t & body );
\end{verbatim}

The definition of the arguments is identical to the standard parallel reduction.
It is important to note, that the deterministic computation does not reproduce sequential computation. As an example, computing

$$\sum_{i=1}^{n} \frac{(-1)^{i+1}}{i}$$

for \( n = 10000000 \) sequentially via

```cpp
for ( auto v : x ) sum += v;
```

gives

6.9314713055990318e-01

If the same reduction is computed with

```cpp
parallel_deterministic_reduce( blocked_range< size_t >( 0, n ),
    double(0),
    [\&x]( const blocked_range< size_t > & r, double s ) { ... }
    std::plus< double >() );
```

the output is

6.9314713055994759e-01

Furthermore, the output is dependent on the grain size.
General Loops

For all previous loops the loop index set was known before parallelisation. Furthermore, a random access was implicitly required to access elements of individual tasks, e.g. vector coefficients to work on.

If this is not the case, parallel_for can not be used for loop parallelisation.

Instead, TBB provides the \textit{parallel_do} algorithm:

\begin{verbatim}
template < typename iterator, typename body_t >
void parallel_do ( iterator first,
                  iterator last,
                  body_t body )
\end{verbatim}

The iterator type for \textit{parallel_do} only has to provide a minimal function set, e.g. construction and increment. Hence, even \texttt{std::list} may be used.

If the iteration space is fixed, e.g. but still restricted to \texttt{std::list}, \texttt{body_t} has to provide only evaluation per set item:

\begin{verbatim}
struct body_t {
    void operator () ( [const] item_t & item ) const; // "const" depends on iterator
};
\end{verbatim}
Fetching new tasks from the iterator in parallel_do is considered a critical section, hence performed strictly sequential.

Therefore, the tasks should be sufficiently large to not let management overhead dominate.

As an example, for a list of double values the following computation was performed on a 2-CPU Xeon E5-2640 with an increasing number of iteration:

```cpp
parallel_do( x.begin(), x.end(), []( double & v ) {
    for ( int it = 0; it < N_ITER; ++it )
        v = std::sin( v );
});
```

The speedup depending on the per task work is shown in the left diagram.

Especially for tasks with little work, parallel_do is not usable for an efficient parallelisation.
General Loops

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The speedup depending on the per task work is show in the left diagram.

Especially for tasks with little work, `parallel_do` is `not` usable for an efficient parallelisation.

The main reason for the inefficiency is the iterator of `std::list`. Using random iterators, the management overhead is much smaller.
Fetching new tasks from the iterator in `parallel_do` is considered a critical section, hence performed strictly sequential.

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Especially for tasks with little work, `parallel_do` is not usable for an efficient parallelisation.

The main reason for the inefficiency is the iterator of `std::list`. Using random iterators, the management overhead is much smaller.
The parallel_do algorithm also permits new tasks to be added to the task set during execution.

For this, the evaluation function of the body_t object has to accept an additional argument:

```cpp
struct body_t {
    void operator () ( [const] item_t & item, 
                      parallel_do_feeder< item_t > & feeder ) const;
};
```

The parallel_do_feeder object provides a function add to extend the task set:

```cpp
parallel_do( x.begin(), x.end(), 
              []( item_t & item, 
                  parallel_do_feeder< item_t > & feeder ) {
                handle( item );

                if ( more_work_available() )
                    feeder.add( more_work() );
            } );
```
Pipeline

For a set of input items, to which functions \( f_1, \ldots, f_m \) are successively applied, e.g. \( f_m(\ldots f_3( f_2( f_1(\cdot))))) \), the pipeline algorithm model may be used:

| input | \( f_1 \) | \( f_2 \) | \( f_3 \) | \( \ldots \) | \( f_m \) | output |

This model is provided by TBB by the \texttt{parallel_pipeline} algorithm, while each stage in the pipeline corresponds to a filter:

```cpp
void
parallel_pipeline ( size_t max_parallel,
                    const filter_t< void, void > & filter_chain );
```

The filter chain itself consists of concatenated filters and can be constructed using \texttt{make_filter}:

```cpp
make_filter< void, outinp_t >( modei, input ) & // feed pipeline
make_filter< input1_t, output1_t >( mode1, func1 ) & // apply \( f_1 \)
make_filter< input2_t, output2_t >( mode2, func2 ) & // apply \( f_2 \)
... make_filter< inputm_t, outputm_t >( modem, funcm ) & // apply \( f_m \)
make_filter< inpout_t, void >( modeo, output ) ); // take items out
```

The output type of filter \( j \) must be identical to the input type of filter \( j + 1 \). The input type of the first and the output type of the last filter have to be \texttt{void}.
The *mode* of a filter may be one of

- `filter::serial_in_order` applies function to items one at a time and always in the same order,
- `filter::serial_out_of_order` applies function to items one at a time in no particular order,
- `filter::parallel` applies function to several items in parallel.

The maximal degree of concurrency for a filter is defined by the value of `max_parallel`. It prevents parallel filters to accumulate an arbitrary number of items, if a following serial filter can not handle the input stream fast enough.

Except for the first function, all function objects `func` have to implement

```cpp
output_t func::operator () ( input_t item );
```

with `output_t` and `input_t` being the input and output types of the corresponding filter.
As the first function object feeds the pipeline, it also controls the end of the pipeline via a `flow_control` object:

```
outinp_t func::operator ( flow_control & fc );
```

If all input items are processed, `fc.stop()` has to be called. The return value of the function is afterwards `not` used:

```
output0_t func::operator ( flow_control & fc ) {
    if ( input.empty() ) {
        fc.stop(); // signal pipeline stop
        return output0_t(0); // return dummy value
    } else {
        return next( input );
    }
}
```
For the introductory setting, the full source using lambda functions is:

```cpp
parallel_pipeline(
  max_parallel,

  // insert items into pipeline
  make_filter< void, double >(
    filter::serial_in_order, [&input] ( flow_control & fc ) {
      if ( input.empty() ) {
        fc.stop();
        return 0.0;
      } else {
        double v = * (input.begin());
        input.pop_front();
        return v;
      }
    }
  )
  &
  // apply f1 ... fm in parallel
  make_filter< double, double >( filter::parallel, [] ( double v ) { return f1( v ); } )
  &
  make_filter< double, double >( filter::parallel, [] ( double v ) { return f2( v ); } )
  &
  make_filter< double, double >( filter::parallel, [] ( double v ) { return f3( v ); } )
  &
  ...
  &
  make_filter< double, double >( filter::parallel, [] ( double v ) { return fm( v ); } )
  &
  // finish pipeline
  make_filter< double, void >(
    filter::serial_in_order, [&output] ( const double v ) {
      output.push_back( v );
    }
  );
)
```

Kriemann, »Introduction to Parallel Programming« 39
Filters may also be created in advance using the `filter_t` class:

```cpp
filter_t< void, double > fin( filter::serial_in_order, 
    [&input] ( flow_control & fc ) { 
        if ( input.empty() ) { 
            fc.stop();
            return 0.0;
        } else { 
            double v = * (input.begin());
            input.pop_front();
            return v;
        } } );

filter_t< double, double > ff1( filter::parallel, [] ( double v ) { return f1( v ); } );
filter_t< double, double > ff2( filter::parallel, [] ( double v ) { return f2( v ); } );
filter_t< double, double > ff3( filter::parallel, [] ( double v ) { return f3( v ); } );
filter_t< double, double > ff4( filter::parallel, [] ( double v ) { return f4( v ); } );
filter_t< double, void > fout( filter::serial_in_order, 
    [&output] ( const double v ) { 
        output.push_back( v );
    } );

filter_t< void, void > f = fin & ff1 & ff2 & ff3 & ff4 & fout; // create pipeline
parallel_pipeline( max_parallel, f );
```

This also allows to `reuse` filters in a pipeline:

```cpp
filter_t< void, void > f = fin & ff1 & ff2 & ff3 & ff4 & ff4 & fout;
```
Non-Linear Pipelines

Pipelines in TBB have to be linear, e.g.

\[ f_1 \rightarrow f_2 \rightarrow f_3 \rightarrow \cdots \rightarrow f_m \]

Multiple inputs or outputs are not supported, e.g.

\[ g_1 \rightarrow g_2 \rightarrow g_3 \rightarrow g_4 \rightarrow g_5 \]

Instead, such pipelines have to be linearised:

\[ g_1 \rightarrow g_2 \rightarrow g_3 \rightarrow g_4 \rightarrow g_5 \]
Pipeline Throughput

The rate at which items are processed by a pipeline depends on several parameters:

- the value of max_parallel,
- slowest filter,
- granularity.

Choosing max_parallel too small may prohibit using all available parallel resources. On the other hand, a value too large may induce too much overhead, e.g. memory.

In any case, the filter with the least throughput rate will determine the throughput of the whole pipeline. This especially applies to serial or IO-bound filter.

Finally, having many filters in the pipeline executing only little work will create more management overhead. Hence, each function should have a reasonable runtime for efficient parallelisation.
Beside the `parallel_pipeline` algorithm, also a `pipeline` class exists in TBB:

```cpp
class pipeline {
public:
    void add_filter ( filter & f ); // add filter to pipeline
    void run ( size_t max_parallel ); // execute pipeline
    void clear(); // remove all filters
};
```

However, in contrast to the `filter_t` class, `filter` is not type safe. New filters have to be derived from `filter` and need to overload the `()` operator:

```cpp
class my_filter : public filter {
    void * operator () ( void * item );
};
```

Input items have to be cast to the actual type, while ensuring that all items have a lifetime equal to the pipeline itself.

The end of the input stream is signaled by a NULL return value of the first filter and the output of the last filter is ignored.
Task Groups
TBB provides a high-level interface to tasks, which allows task definition by using standard C++ function objects.

Each of these tasks belongs to a *task group* implemented by the class `task_group`:

```cpp
class task_group {
    template< typename func_t > void run ( const func_t & f );
    task_group_status void wait ();
    template< typename func_t > void run_and_wait ( const func_t & f );
};
```

Spawning a task in a task group is performed using the function `run()`, whereas `wait()` will block until all tasks in the task group have finished execution:

```cpp
task_group g;
g.run( [] { ... } );
g.run( [] { ... } );
g.wait();
```

Spawning the last task in a group and waiting for all tasks is combined in `run_and_wait()`.
Task Groups

Beside spawning and waiting, tasks in a task group may be cancelled via

```cpp
void task_group::cancel ();
```

Tasks, which are not yet running will *not* be scheduled for execution if the group has been cancelled. However, if a task is already running, cancellation has *no* effect.

The return value of the function `wait()` gives information about the completion and cancellation status of the task group using the `task_group_status` enum:

- `not_complete`: group was not cancelled and there are still tasks which have not completed,
- `complete`: group was not cancelled and all tasks completed,
- `canceled`: group was cancelled.

The current cancellation status can also be retrieved using the function

```cpp
void task_group::is_cancelling ();
```

**Remark**

*If an exception is thrown within a task, the task group is cancelled.*
Example: Dot Product

For computing the dot product \( \langle x, y \rangle \) with \( x, y \in \mathbb{R}^n \), a recursive algorithm is chosen. Each recursive call corresponds to a new task, belonging to a new task group:

```cpp
double parallel_dot ( const std::vector< double > & x,
 const std::vector< double > & y,
 const size_t lb, const size_t ub ) {
    if ( ub - lb < MIN_SIZE ) {
        return std::inner_product( & x[lb], & x[ub], & y[lb], 0.0 );
    } else {
        task_group g;
        double res1 = 0;
        double res2 = 0;

        g.run( [ub,lb,&res1,&x,&y] { res1 = parallel_dot( x, y, lb, (ub-lb)/2 ); } );
        g.run( [ub,lb,&res2,&x,&y] { res2 = parallel_dot( x, y, (ub-lb)/2, ub ); } );
        g.wait();

        return res1 + res2;
    }
}
```

Kriemann, »Introduction to Parallel Programming« 47
Tasks
The task Class

The fundamental building block of the low-level task based programming in TBB is the `task` class and its `execute` function:

```cpp
class task {
public:
    virtual task * execute() = 0;
    ...
};
```

All user tasks have to be derived from `task` and overload `execute`.

Each task has an optional `parent` and a `reference counter` associated with it, accessible via

```cpp
    task * task::parent () const;
    void task::set_parent ( task * p );
    void task::set_ref_count ( int count );
    void task::increment_ref_count ();
    int task::decrement_ref_count ();
```

The reference counter of a task \( t \) contains the number of tasks to which \( t \) is the parent. Modifications of a reference counter are always atomic.

A `root` task has no parent and a reference counter of 0. Furthermore, if the `parent` field of a task is set, this task is a `child` task of the corresponding parent.
Tasks in TBB provide special allocation routines for an efficient management of tasks.

```cpp
static proxy1_t allocate_root ();
proxy2_t allocate_child ();
static proxy3_t allocate_additional_child_of ( task & t );
```

**Remark**

The allocation function return internal proxy objects, which handle the actual allocation.

Depending on the allocation routine, the parent and the reference counter of the new (or calling) task is set:

<table>
<thead>
<tr>
<th>Allocation Routine</th>
<th>Parent</th>
<th>Reference Counter</th>
</tr>
</thead>
<tbody>
<tr>
<td>allocate_root()</td>
<td>nullptr</td>
<td>0</td>
</tr>
<tr>
<td>t-&gt;allocate_child()</td>
<td>t</td>
<td>0</td>
</tr>
<tr>
<td>allocate_additional_child_of( t )</td>
<td>t</td>
<td>0, t.refcount++</td>
</tr>
</tbody>
</table>

*All* tasks must be allocated by these functions. Otherwise, the results of the allocation is undefined.
TBB provides several methods to spawn new tasks and synchronise with the end of tasks, depending on how the tasks were allocated.

For a parent task \( p \), a child task \( t \) is put into the thread-local work set via

```c++
static void task::spawn ( task & t );
```

For further synchronisation to work, the value of the reference counter of \( p \) is important. Before spawning any task, it has to be set to the number of child tasks.

After spawning, \( p \) can wait for all child tasks to finish by calling

```c++
void task::wait_for_all ();
```

For this function, an additional reference is needed in \( p \), because \( \text{wait} \_\text{for} \_\text{all}() \) will execute tasks in the work set, until the reference counter of \( p \) is 1. Afterwards, it decreases the reference counter of \( p \) to 0 and returns.

The reference counter of \( p \) is decremented by using the parent variable in child tasks after \( \text{execute}() \) has finished. If it becomes 0, the parent is also put in the work set by TBB.
Spawning the last and waiting for all child tasks can be combined by

```cpp
void task::spawn_and_wait_for_all ( task & t );
```

**Remark**

_The task spawned by `spawn_and_wait_for_all()` is guaranteed to be executed by the current thread._

For root tasks a special spawn routine is provided:

```cpp
static void task::spawn_root_and_wait ( task & t );
```

**Destruction**

Tasks in the work set are usually automatically destroyed upon finishing execution. The task memory is afterwards recycled for future allocations.

If this is not the case, e.g. due to special handling of the reference counter, tasks should be explicitly destroyed using

```cpp
static void task::destroy ( task & t );
```

The reference counter of `t` must be 0. If a parent of `t` exists, its reference counter is decreased.
Compute \( \langle x, y \rangle \) recursively, while each recursive call forms a new task:

```cpp
class dot_task_t : public task {
public:
    const std::vector<double> & x, & y;
    const size_t lb, ub; // interval [lb,ub)
    double & res; // holds locally computed result

task * execute () {
    if ( ub - lb < MIN_SIZE ) {
        res = std::inner_product( & x[lb], & x[ub], & y[lb], 0.0 );
    } else {
        double res1 = 0;
        double res2 = 0;
        dot_task_t & child1 = * new ( allocate_child() ) dot_task_t( x, y, lb, (ub+lb)/2, res1 );
        dot_task_t & child2 = * new ( allocate_child() ) dot_task_t( x, y, (ub+lb)/2, ub, res2 );

        set_ref_count( 3 );
        spawn( child2 );
        spawn_and_wait_for_all( child1 );

        res = res1 + res2;
        return nullptr;
    }
} } ;

double parallel_dot ( const std::vector<double> & x, const std::vector<double> & y ) {
    double res = 0;
    dot_task_t & root = * new ( task::allocate_root() ) dot_task_t( x, y, 0, x.size(), res );

    task::spawn_root_and_wait( root );
    return res;
}
```

On an Intel XeonPhi 5110P, the speedup of `parallel_dot()` for \( n = 10^8 \) is 12.63x compared to 2.56x for `parallel_reduce()`.
The task graph for the dot-product example with $n = 8$ and $\text{MIN\_SIZE} = 1$ is:

The computation starts with the $\sum_0^7$ task.

The task scheduler of TBB uses a mixed DFS/BFS scheduling strategy to optimise cache locality and parallel execution:

- tasks per thread are executed via DFS (preserves cache locality),
- task stealing follows a BFS strategy (increases parallel degree)

Some methods are provided by TBB to influence or optimise task scheduling.
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- task stealing follows a BFS strategy (increases parallel degree)

Some methods are provided by TBB to influence or optimise task scheduling.
A thread, which executes a task calling `spawn_and_wait_for_all()`, will executes other tasks while the child tasks are running. If the parent task may finally proceed, the corresponding thread could be busy, preventing immediate execution of the parent task.

This can be overcome by using a `continuation` task, i.e. a separate task which continues execution after all child tasks of the original parent task have finished:

```cpp
task * dot_task_t::execute () {  
    if ( ub - lb < MIN_SIZE ) { ... }  
    else {  
        cont_t & cont = * new ( allocate_continuation() ) cont_t( res );  
        task & child1 = * new ( cont.allocate_child() ) dot_task_t( x, y, lb, (ub+lb)/2, cont.res1 );  
        task & child2 = * new ( cont.allocate_child() ) dot_task_t( x, y, (ub+lb)/2, ub, cont.res2 );  
        cont.set_ref_count( 2 );  
        spawn( child2 );  
        spawn( child1 );  
        return nullptr;  
    } }
```

The child tasks will have the continuation task as their parent. As its reference counter is 2, it is scheduled immediately after the last child finishes. If the current thread is busy, another thread may steal the continuation task and execute it.

**Remark**

*Using a continuation task for the dot-product, the speedup increases to 23.93x.*
Up to now, the return value of `execute()` was always `nullptr`.

It can also be a pointer to a task, which is then chosen as the next task to be executed by the current thread.

```cpp
task * dot_task_t::execute () {
    if ( ub - lb < MIN_SIZE ) { ... } 
    else {
        cont_t & cont = * new ( allocate_continuation() ) cont_t( res );
        task & child1 = * new ( cont.allocate_child() ) dot_task_t( x, y, lb, (ub+lb)/2, cont.res1 );
        task & child2 = * new ( cont.allocate_child() ) dot_task_t( x, y, (ub+lb)/2, ub, cont.res2 );

        cont.set_ref_count( 2 );
        spawn( child2 );

        return & child1;
    }
}
```

Before, the `child1` task was put in the work set of the current thread, and taken out after `execute()` has finished.

Both operations are eliminated by handing `child1` directly to the task scheduler.

**Remark**

*For the dot-product computation, the speedup increases to 29.03x, when bypassing the scheduler.*
Bypassing the scheduling process can be combined with bypassing task allocation and deallocation.

```cpp
task * dot_task_t::execute () {
    if ( ub - lb < MIN_SIZE ) { ... }
    else {
        cont_t & cont = * new ( allocate_continuation() ) cont_t( res );
        task & child2 = * new ( cont.allocate_child() ) dot_task_t( x, y, (ub+lb)/2, ub, cont.res2 );

        recycle_as_child_of( cont );
        ub = (ub+lb)/2;
        res = cont.res1;

        cont.set_ref_count( 2 );
        spawn( child2 );

        return this;
    }
}
```

The current task $t$ is updated with the data of child1. Furthermore,

```cpp```
void task::recycle_as_child_of( task & new_parent );
```

will set the continuation task of $t$ as its new parent.

Finally, the current task is reexecuted immediately by returning a pointer to it.

**Remark**

*If the parent task is recycled as child during dot-product computation, the speedup increases to 31.70x.*
Recycling a parent task as a child task is especially efficient, if the child can reuse data from the parent.

If on the other hand, a continuation task may reuse data, a task can also be recycled as its own continuation via

```cpp
void task::recycle_as_continuation ();
```

A potential race condition exists, if all child tasks of a recycled task \( t \) have finished, at which point \( t \) is scheduled to be executed again.

If \( t \) has not finished its own `execute()` function, it is run in parallel. As this may lead to various side effects, TBB requires that such situations must not occur.

One way to prevent this execution overlap is to *not* spawn one of the child tasks but return it from `execute()`.

Another way is to set the reference count of the continuation task to \( k + 1 \) where \( k \) is the number of child tasks and use

```cpp
void task::recycle_as_safe_continuation ();
```

which avoids the race condition.
Shared Queue

Each worker thread in TBB has its local work set of tasks. If not using scheduler bypass or when the last child of a parent task finished, the thread will look first at this set for a new task to execute.

TBB also has a *shared set* for all threads, from which it will take tasks if the thread local set is empty. Only if this shared set is also empty, it will try to steal tasks from other sets.

New tasks may be explicitly put into the shared task set via

```cpp
static void task::enqueue ( task & t );
```

In contrast to the thread local work set, which is handled in a *last-in first-out* order, the shared set is accessed (roughly) in *first-in first-out* order.

This behaviour for the shared set ensures some *fairness* for task execution, i.e. eventually the enqueued task will be executed. Tasks in the thread-local set may not be scheduled until another task explicitly waits for them.

Furthermore, even if the number of worker threads is zero, a thread is started to handle tasks in the shared set.
Thread affinity and, assuming a processor bound thread, also cache affinity for tasks is supported by TBB through a set of task functions.

The affinity id of an executed task is automatically send to the task via

```cpp
virtual void task::note_affinity ( affinity_id id );
```

By overloading this function in a user defined task, the affinity id can be recorded and set for a later spawned task using

```cpp
void task::set_affinity ( affinity_id id );
```

The task scheduler will afterwards use the thread with the same id to execute the task. For this, the affinity id has to be set `before` the task is spawned:

```cpp
task & child = * new( allocate_child() ) child_task_t( ... );
child.set_affinity( my_id );
spawn( child );
```

Access to the current affinity id of a task is provided with

```cpp
affinity_id task::affinity () const;
```

**Remark**

Variables of type `affinity_id` must default to 0 or hold values from `note_affinity()`.
For tasks in the shared queue a priority may be set as an optional argument to enqueue:

```c++
static void task::enqueue ( task & t, priority_t p );
```

TBB supports three different levels of priority:

- `priority_high`,
- `priority_normal`,
- `priority_low`.

Tasks of a higher priority will be first taken out of the shared queue and scheduled for execution.

The priority of the task is fixed when calling `enqueue()`, i.e. can not be changed afterwards.

**Remark**

*Spawned tasks will have the priority of their task group (see below).*
All tasks are part of a task group, represented by objects of type `task_group_context`.

Each creation of an `task_scheduler_init`, e.g. during explicit TBB initialisation, also creates an implicit task group object. All tasks spawned during the lifetime of that object, automatically belong to this task group:

```c
{
    task_scheduler_init tsi1( 4 ); // task group 1
    spawn_root_and_wait( ... );

    {
        task_scheduler_init tsi2( 8 ); // task group 2
        spawn_root_and_wait( ... );
    }
}
{
    task_scheduler_init tsi3( 2 ); // task group 3
    spawn_root_and_wait( ... );
}
```

All task groups in a process form a forest of trees, i.e. multiple root task groups may exist, each having various sub groups.

The hierarchy defined by the tree also determines certain properties of the task groups, e.g. for cancellation or priority.
A root task may also be explicitly assigned to a task group by providing the `task_group_context` object during task allocation:

```cpp
    task_group_context g;
    task & t = * new( task::allocate_root( g ) ) my_task_t( ... );
```

All child tasks allocated with `task::allocate_child()` will automatically belong to the same task group as the parent task.

Changing the task group of a task can be done using

```cpp
    void task::change_group ( task_group_context & g );
```

The task group of a task is accessed via

```cpp
    task_group_context * task::group ();
```
All tasks in a task group and of all sub groups can be cancelled using

```cpp
bool task::cancel_group_execution();
```

Tasks already running are not directly affected by this. However, they may query the current cancellation status via

```cpp
bool task::is_cancelled();
```

and stop further execution, e.g.:

```cpp
task * execute () {
    while ( ! is_cancelled() ) {
        proceed_computation();
    }
}
```

All other tasks, i.e. tasks still waiting for execution, will not call `execute()` after a cancellation request was issued.

If the task group already received a cancellation request, the return value of `cancel_group_execution()` is false.

**Remark**

*Throwing an exception within a task will also cancel further execution within the corresponding task group.*
Associated with a task group is a *priority* level of the same type and values as for tasks in the shared queue, i.e. either priority_high, priority_normal or priority_low.

The priority of a task group affects all tasks in the group and in all sub groups. Tasks in a task group of a higher priority will be executed before tasks in group with a lower priority.

To set the priority of a task group either use

```cpp
void task_group_context::set_priority ( priority_t p );
```

or

```cpp
void task::set_group_priority ( priority_t p );
```

Similarly, the priority of a task group may be accessed using one of

```cpp
priority_t task_group_context::priority ();
priority_t task::group_priority ();
```

In contrast to the priority of an enqueue’d task, the priority of task groups may be changed at any time. However the effect may not be immediate, e.g. tasks with a lower priority may still be scheduled first.
Spawning several tasks can also be combined using the TBB type `task_list`:

```c++
static void spawn ( task_list & list );
void spawn_and_wait_for_all ( task_list & list );
static void spawn_root_and_wait ( task_list & list );
```

The class `task_list` supports basic STL container functions:

```c++
class task_list {
public:
    bool empty () const; // return true if empty
    void push_back ( task & task ); // insert task at end of list
    task & pop_front (); // remove and return first task in list
    void clear (); // remove all tasks from list
};
```

Spawning task lists may be more efficient than spawning each task separately:

```c++
task_list list;
set_ref_count( nchildren + 1 );
for ( int i = 0; i < nchildren; ++i )
    list.push_back( * new( allocate_child() ) child_task_t( ... ) );
spawn_and_wait_for_all( list );
```

**Remark**

*Recursive task spawning should be preferred over long task lists.*
Up to now, task creation was coupled with task spawning due to recursive parallelism, which also reflected in the relation between parent and child tasks.

The task concept can also be viewed in terms of *successors* and *predecessors*, where a successor task is spawned after a predecessor task has finished.

This relation can be expressed in terms of a graph, where nodes correspond to tasks and directed edges to dependencies between them. Since no cyclic dependencies may exist, such tasks are *directed acyclic graphs* or *DAGs*.

For the block-wise LU factorisation of a dense matrix $A \in \mathbb{R}^{n \times n}$ the incoming and outgoing dependencies for a single matrix block are shown to the right.

Incoming dependencies are due to needed matrix updates (blue) and the diagonal factorisation (red). Outgoing dependencies are because of matrix updates to blocks below (green).
Since several incoming dependencies exist, pre-allocation of tasks is of advantage. Here, the reference counter may be set according to the degree of incoming edges in the DAG. After each predecessor has finished, it can decrement the reference counter of all successors. If this reaches zero, the task is scheduled for execution.

For the block-wise LU factorisation, matrix block $A_{ij}$ has $2 \cdot \min(i, j)$ dependencies due to matrix updates. In addition, off-diagonal blocks depend on the factorisation of the diagonal block.

This leads to the following implementation, where for each matrix block a task is pre-allocated, either for factorisation or matrix solves, and initialised with the corresponding number of dependencies:

```cpp
std::vector< std::vector< task * > > tasks;

for ( size_t i = 0; i < b; ++i ) {
    // factorise diagonal block
    tasks[i][i] = new( task::allocate_root() ) fac_task_t( A, i, tasks, mutexes );
    tasks[i][i]->set_ref_count( 2*i );

    // solve offdiagonal blocks
    for ( size_t j = i+1; j < b; ++j ) {
        tasks[i][j] = new( task::allocate_root() ) solve_lower_task_t( A, i, j, tasks, mutexes );
        tasks[i][j]->set_ref_count( 1 + 2*i );

        tasks[j][i] = new( task::allocate_root() ) solve_upper_task_t( A, j, i, tasks, mutexes );
        tasks[j][i]->set_ref_count( 1 + 2*i );
    }
}
```
The implementation of a factorisation task consists of the computation of all updates, the actual point-wise factorisation and spawning successor tasks:

```cpp
class fac_task_t : public task {
private:
    task_matrix_t & tasks;

public:
    task * execute () {
        const size_t b = A.n/BLOCK_SIZE;
        Matrix A_ii( BLOCK_SIZE );

        // spawn update tasks
        spawn_updates( A, this, i, i, i );

        // factorise
        load_block( A, A_ii, i, i );
        lu( A_ii );
        store_block( A, A_ii, i, i );

        // spawn tasks below and to the right
        for ( size_t k = i+1; k < b; ++k ) {
            if ( tasks[i][k]->decrement_ref_count() == 0 ) spawn( * tasks[i][k] );
            if ( tasks[k][i]->decrement_ref_count() == 0 ) spawn( * tasks[k][i] );
        }

        return nullptr;
    }
};
```

Tasks for matrix solves are similar, e.g. first apply updates, then perform matrix solve and finally spawn successors.
The function `spawn_updates()` spawns tasks for all matrix updates

\[
A_{ij} = A_{ij} - \sum_{k=0}^{\ell-1} L_{ik} U_{kj}
\]

with \( \ell = \min(i, j) \).

In contrast to factorisation and solve tasks, update tasks are created and spawned when needed:

```c
void spawn_updates ( Matrix & A, task * t size_t i, size_t j, size_t l ) {
    t->set_ref_count( l+1 );
    // spawn updates for
    // A_{ij} = A_{ij} - \sum_{k=0}^{\ell-1} L_{ik} U_{kj}
    for ( size_t k = 0; k < l; ++k )
        t->spawn( * new( t->allocate_child() )
            update_task_t( A, i, j, k ) );
    t->wait_for_all();
}

class update_task_t : public task {
... task * execute () {
    Matrix U_sub( BLOCK_SIZE );
    Matrix L_sub( BLOCK_SIZE );
    Matrix A_sub( BLOCK_SIZE );
    load_block( A, L_sub, i, k );
    load_block( A, U_sub, k, j );
    load_block( A, A_sub, i, j );
    multiply_sub( L_sub, U_sub, A_sub );
    return nullptr;
} }
```
DAG Computations

The function `spawn_updates()` spawns tasks for all matrix updates

\[ A_{ij} = A_{ij} - \sum_{k=0}^{\ell-1} L_{ik} U_{kj} \]

with \( \ell = \min(i, j) \).

In contrast to factorisation and solve tasks, update tasks are created and spawned when needed:

```cpp
void spawn_updates ( Matrix & A, task * t
size_t i, size_t j, size_t l ) {
    t->set_ref_count( l+1 );
    // spawn updates for
    A_{ij} = A_{ij} - \sum_{0 \leq k < l} L_{ik} U_{kj}
    for ( size_t k = 0; k < l; ++k )
        t->spawn( * new( t->allocate_child() )
            update_task_t( A, i, j, k ) );
    t->wait_for_all();
}
```

```cpp
class update_task_t : public task {
...
task * execute () {
    Matrix U_sub( BLOCK_SIZE );
    Matrix L_sub( BLOCK_SIZE );
    Matrix A_sub( BLOCK_SIZE );
    load_block( A, L_sub, i, k );
    load_block( A, U_sub, k, j );
    load_block( A, A_sub, i, j );
    multiply_sub( L_sub, U_sub, A_sub );
    return nullptr;
}
```

**Remark**

*Updating a matrix block represents a critical section (see below).*
DAG Computations

The node to start the whole computation corresponds to the matrix block $A_{00}$. The end of the process is formed by the task for matrix block $A_{b-1,b-1}$.

Furthermore, by construction, a path exists from the start to the final node in the DAG.

The execution of the DAG is initiated by the final node, with the start node as the first task to execute.

As before, an additional reference is needed for `spawn_and_wait_for_all()`. Since this prevents automatic execution of the final node, it is performed manually, as is the destruction of the task:

```cpp
tasks[b-1][b-1]->increment_ref_count();
tasks[b-1][b-1]->spawn_and_wait_for_all( * tasks[0][0] );
tasks[b-1][b-1]->execute();
task::destroy( * tasks[b-1][b-1] );
```
DAG Computations

The node to start the whole computation corresponds to the matrix block $A_{00}$. The end of the process is formed by the task for matrix block $A_{b-1,b-1}$.

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tasks[b-1][b-1]->increment_ref_count();
tasks[b-1][b-1]->spawn_and_wait_for_all( *tasks[0][0] );
tasks[b-1][b-1]->execute();
task::destroy( *tasks[b-1][b-1] );
```
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tasks[b-1][b-1]->increment_ref_count();
tasks[b-1][b-1]->spawn_and_wait_for_all( *tasks[0][0] );
tasks[b-1][b-1]->execute();
task::destroy( *tasks[b-1][b-1] );
```
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The execution of the DAG is initiated by the final node, with the start node as the first task to execute.

As before, an additional reference is needed for `spawn_and_wait_for_all()`. Since this prevents automatic execution of the final node, it is performed manually, as is the destruction of the task:

```cpp
tasks[b-1][b-1]->increment_ref_count();
tasks[b-1][b-1]->spawn_and_wait_for_all( *tasks[0][0] );
tasks[b-1][b-1]->execute();
task::destroy( *tasks[b-1][b-1] );
```

The parallel speedup for $n = 8192$ and a block size of 64 is:

<table>
<thead>
<tr>
<th></th>
<th>Xeon X5650 (24 threads)</th>
<th>Xeon E5-2640 (24 threads)</th>
<th>XeonPhi 5110P (240 threads)</th>
</tr>
</thead>
<tbody>
<tr>
<td>DAG-based</td>
<td>9.11x</td>
<td>9.94x</td>
<td>72.36x</td>
</tr>
</tbody>
</table>
Mutual Exclusion
Mutual Exclusion

TBB implements various forms of mutexes. The differences between these mutex classes are based on certain properties mutexes may have:

**Scalable:** Using a *scalable* mutex will not result in a worse performance than sequential execution, e.g. due to high processor or memory bandwidth usage.

**Fair:** A *fair* mutex guarantees, that all threads waiting for the mutex, will eventually succeed in locking it.

**Recursive:** A *recursive* mutex permits a thread, who already holds a lock on a mutex to lock it again.

**Yield or Block:** A *yielding* mutex will only temporarily yield a processor to other threads, while a *blocking* mutex will yield the processor until the lock is acquired.

**Remark**

Unscalable mutexes may be faster if the lock contention is low. Similarly, fairness requires more overhead than unfair behaviour. Reactivating a blocking thread also takes some time. Hence, yielding should be preferred if the waiting time for a mutex is short.
In the following table, the different mutex types of TBB with their properties are shown:

<table>
<thead>
<tr>
<th>Mutex Type</th>
<th>Scalable</th>
<th>Fair</th>
<th>Recursive</th>
<th>Long Wait</th>
</tr>
</thead>
<tbody>
<tr>
<td>mutex</td>
<td>OS dep.</td>
<td>OS dep.</td>
<td>no</td>
<td>blocks</td>
</tr>
<tr>
<td>recursive_mutex</td>
<td>OS dep.</td>
<td>OS dep.</td>
<td>yes</td>
<td>blocks</td>
</tr>
<tr>
<td>spin_mutex</td>
<td>no</td>
<td>no</td>
<td>no</td>
<td>yields</td>
</tr>
<tr>
<td>queuing_mutex</td>
<td>yes</td>
<td>yes</td>
<td>no</td>
<td>yields</td>
</tr>
<tr>
<td>null_mutex</td>
<td>/</td>
<td>yes</td>
<td>yes</td>
<td>never</td>
</tr>
</tbody>
</table>

Both, `mutex` and `recursive_mutex` are based on the implementation of mutexes of the operation system.

All other mutexes are implemented in TBB, i.e. no operating system functionality is used.

A special case is `null_mutex`, which can be used as a dummy.
Scoped Locking

Locking and unlocking mutexes manually easily results in program errors, e.g. due to forgotten mutex release. Also, unlocking mutexes in case of an exception is tedious and error-prone.

Instead, mutexes should only accessed via scoped locks. Each mutex type $M$ in TBB has an associated type for a scoped lock:

\[
M::\text{scoped\_lock}
\]

with the following interface:

```cpp
class M::scoped_lock {
    scoped_lock (); // default constructor without mutex
    scoped_lock ( M & m ); // constructor: lock mutex m
    ~scoped_lock (); // unlock mutex if previously locked
    void acquire ( M & m ); // lock mutex m
    bool try_acquire ( M & m ); // try to lock mutex m; return true if successful
    void release (); // unlock previously locked mutex
};
```

By definition, unlocking is coupled with the destruction of the scoped lock, and hence bound to the object lifetime.

Therefore, scoped locks should be auto variables, e.g. not allocated via new.
In a typical scenario, the scoped lock is constructed at the beginning of a critical section and automatically destroyed at its end:

```cpp
class spin_mutex { /* ... */

spin_mutex m;
...
{
    spin_mutex::scoped_lock lock( m ); // acquire lock
    ...
} // release lock automatically
...
```

or

```cpp
void f ( mydata_t & d, recursive_mutex & m ) {
    recursive_mutex::scoped_lock lock( m ); // acquire lock
    ...
} // release lock automatically
```

### C++11 Interface

TBB mutexes also support the standard C++11 mutex functions:

```cpp
void M::lock () ; // lock mutex
bool M::try_lock () ; // try to lock mutex
void M::unlock () ; // unlock mutex
```

This also allows to use C++11 std::lock_guard together with TBB mutexes.
Reader/Writer Locks

TBB extends mutexes to distinguish between *readers* and *writers* with respect to a critical section. Multiple readers may hold a lock, but only a single writer.

The extension affects only the special mutex classes from TBB, not the OS related mutexes.

For each TBB mutex, a corresponding Reader/Writer (RW) lock exists:

<table>
<thead>
<tr>
<th></th>
<th>Scalable</th>
<th>Fair</th>
<th>Recursive</th>
<th>Long Wait</th>
</tr>
</thead>
<tbody>
<tr>
<td>spin_rw_mutex</td>
<td>no</td>
<td>no</td>
<td>no</td>
<td>yields</td>
</tr>
<tr>
<td>queuing_rw_mutex</td>
<td>yes</td>
<td>yes</td>
<td>no</td>
<td>yields</td>
</tr>
<tr>
<td>null_rw_mutex</td>
<td>/</td>
<td>yes</td>
<td>yes</td>
<td>never</td>
</tr>
</tbody>
</table>

The scoped lock interface is also extended, now accepting a boolean parameter `writer` with each mutex, indicating whether the lock shall be a write lock (`true`) or a reader lock (`false`):

```cpp
class M::scoped_lock {
    scoped_lock ( M & m, bool writer );
    void acquire ( M & m, bool writer );
    bool try_acquire ( M & m, bool writer );
};
```
Reader/Writer Locks

Reader locks may also be upgraded to writer locks and vice versa:

```cpp
bool M::scoped_lock::upgrade_to_writer (); // upgrade to write lock
bool M::scoped_lock::downgrade_to_reader (); // downgrade to reader lock
```

In both cases, either a reader (for `upgrade_to_writer()`) or a writer lock (for `downgrade_to_reader()`) must already be acquired. Otherwise, the effect is undefined.

When upgrading a reader lock to a writer lock, it will block until the last reader releases the lock.

**Remark**

*Both functions return false if the lock was released and reacquired, e.g. via* `release()`/*acquire()*, *and true otherwise.*

An example usage of RW locks are read-only and write sections in a program:

```cpp
{
    M::scoped_lock lock( m, false ); // acquire reader lock
    // access data read-only
    {
        lock.upgrade_to_writer();
        // write data
        lock.downgrade_to_reader();
    }
    // access read-only again
    ...
}
```
Containers
Containers

Standard C++ or C++11 STL containers are *not* multi-thread safe if modified. Hence, access has to be guarded by a mutex, preventing efficient parallel operations.

TBB on the other hand, implements several containers, which permit concurrent access, e.g. insertion, deletion, by using either

- a *fine-grained* locking mechanism, i.e. locking only small parts of the data structure which are actually changed or
- *lock-free* mechanisms, where threads notice changes by other threads and handle the effects of these changes automatically.

The implemented containers fall into the following categories:

- **dynamic arrays:** implementes dynamic resize of an array
- **associative arrays:** a *key* type is used to access data,
- **queues:** implements *first-in first-out* data access.

**Remark**

*Concurrent access comes with a cost, i.e. worse sequential performance. Therefore, concurrent containers should be used, if the advantage of concurrency and the corresponding speedup outweighs the additional costs.*
The concurrent version of std::vector in TBB is **concurrent_vector**.

It supports concurrent growing of the vector, i.e. several threads increase the size simultaneously.

For this, the following methods are provided:

```cpp
iterator concurrent_vector::grow_by(size_t delta [, const item_t & t ]); 
iterator concurrent_vector::grow_to_at_least(size_t n ); 
iterator concurrent_vector::push_back(const item_t & t );
```

Here, `item_t` corresponds to the template argument of `concurrent_vector`.

The function `grow_by()` will append `delta` new default constructed items (or copies of `t`) to the array. With `grow_to_at_least()`, the new array size is at least `n`. Finally, `push_back()` will append `t` to the end of the vector.

**Remark**

*Consecutive items in a concurrent vector may not necessarily have a consecutive address!*
When growing a concurrent array, access to array items is still possible in parallel. However, new items may *not* yet have been constructed.

Therefore, when fetching items from an array, further tests are necessary to ensure correctness when accessing item $i$:

- ensure, that the array size is at least $i$,
- ensure, that the element was constructed.

While the first test is simple, e.g. by using the `size()` function, to ensure that an object was constructed, is more involved.

For this, an *atomic* flag may be used as part of the object class, which is initialised last during construction. Also the value of the (atomic) element itself may signal construction, if it differs from some default value, e.g. 0.

The general procedure may look like:

```cpp
item_t fetch_item ( const concurrent_vector< item_t > & v, const size_t i ) {
    while ( v.size() <= i ) // test vector size
        std::this_thread::yield();
    while ( ! v[i].is_constructed ) // test object construction
        std::this_thread::yield();
    return v[i];
}
```
An associative array, also called *hash map*, uses a *key* to access data. For this, a hash value of the key is computed and used as an index to some internal data structure (usually an array).

C++11 provides associative arrays as `unordered_map` and `unordered_multimap`. TBB extends these for concurrent access as `concurrent_unordered_map` and `concurrent_unordered_multimap`.

Both container types support concurrent insertion and traversal, but *not* concurrent deletion of elements.

The following functions are safe for concurrent access:

<table>
<thead>
<tr>
<th>Function</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><code>bool empty()</code></td>
<td>Checks if the map is empty.</td>
</tr>
<tr>
<td><code>size_t size()</code></td>
<td>Returns the size of the map.</td>
</tr>
<tr>
<td><code>pair&lt;iterator, bool&gt; insert(const value_t &amp; x)</code></td>
<td>Inserts the element into the map.</td>
</tr>
<tr>
<td><code>size_t count(const key_t &amp; k)</code></td>
<td>Returns the count of elements with the given key.</td>
</tr>
<tr>
<td><code>pair&lt;iterator, iterator&gt; equal_range(const key_t &amp; k)</code></td>
<td>Returns an iterator range that includes all elements with the given key.</td>
</tr>
</tbody>
</table>

Furthermore, `concurrent_unordered_map` provides a multi-thread safe index operator `[]`. 
Associative Arrays

Furthermore, safe iteration via `begin()` and `end()` is possible. The corresponding iterators remain valid, even if a new element is inserted.

All of these methods are *lock-free*, i.e. no user visible mutex is used.

TBB also implements `concurrent_hash_map`, which furthermore provides concurrent deletion, albeit with locks. Beside that, the interface is identical to the above classes.

An important difference affects iteration. It is *not* safe to use concurrent methods while iterating over a `concurrent_hash_map`.

**Hashing**

Hash values of keys are computed in TBB via the `tbb_hash` class

```cpp
class tbb_hash {
    size_t operator () ( const key_t & k ) { return tbb_hasher( k ); } 
};
```

For user defined keys, a new function `tbb_hasher()`

```cpp
size_t tbb_hasher ( const key_t & k );
```

may be implemented to use the above container classes.
Similar to associative arrays, C++11 defines associative sets which store elements accessed by a key value with no particular order, thereby enabling faster access.

Again, TBB extends these with concurrent insertion and traversal as `concurrent_unordered_set` and `concurrent_unordered_multiset`.

The interface is identical with the associative array counterparts, with the exception of an index operator.

Furthermore, hashing is identical to associative arrays, i.e. via `tbb_hasher` function.
The concurrent version of the C++ queue class in TBB is `concurrent_queue`.

It provides the basic functionality of a queue: appending items at its end and removing items from the top of the queue.

However, if multiple threads access a shared queue simultaneously, items at the top of the queue may no longer be available when trying to remove them.

Therefore, in TBB the `pop_front()` function is enhanced by a test for an item available. The resulting function is `try_pop()`:

```
C++ Sequential Access  Parallel Access

std::queue< int > q;
... while ( ! q.empty() ) {
    int i = q.pop_front();
    ...
}

concurrent_queue< int > q;
... int i;
while ( q.try_pop( i ) ) {
    ...
}
```

The return value of `try_pop()` is `true`, if an item was available, which is assigned to the function argument.

The counterpart of `try_pop()` is `push()`, which is equivalent to `push_back()` from `std::queue`. 

Kriemann, »Introduction to Parallel Programming« 86
Queues

Beside `push()` and `try_pop()`, `concurrent_queue` also has several functions equivalent to `std::queue` functions. But since these functions are not safe in a multi-threaded environment, they are prefixed with `unsafe`:

```cpp
size_t concurrent_queue::unsafe_size  () const;
iterator concurrent_queue::unsafe_begin ();
iterator concurrent_queue::unsafe_end   ();
```

While `concurrent_queue` will handle arbitrary many items, TBB also implements a `bounded` queue: `concurrent_bounded_queue`. Here, a maximal capacity may be specified, which is unlimited by default:

```cpp
size_t concurrent_bounded_queue::capacity  () const;
void    concurrent_bounded_queue::set_capacity ( size_t capacity );
```

As a consequence, pushing items to the queue via `push()` may block until there is enough capacity to insert the new item. With `try_push()`, the push operation is skipped if it would exceed the queue capacity.

Furthermore, `concurrent_bounded_queue` implements a blocking `pop()` function, i.e. it will wait until the queue is non-empty. However, the function `try_pop()` is also available.

To wake up any thread waiting in `push()` or `pop()`, the `abort()` function can be used.
An Argument Against Queues

Since queues provide a FIFO access to data, insertion and removal has to follow a strict sequential order.

**Remark**

*Due to concurrent insertion or removal, TBB only guarantees that if a single thread inserts two items and if another thread removes them, the order is preserved.*

This sequential order may pose a bottleneck of many items are managed by a queue in short time.

Furthermore, due to the FIFO order, the time between data insertion and removal may be long, such that the data may have been removed from the local cache or, a remote thread may have pop’ed the data.

Therefore, it is suggested to use other functions instead of queues, e.g. `parallel_do` (with local insertion) or pipelines. Here, cache and thread locality is preserved.
Miscellanea
TBB implements various other algorithms, some of which are now available in C++11, e.g. atomics or timing.

Others are abbreviations of more general algorithms, e.g. `parallel_invoke()`, or implement standard algorithms, e.g. `parallel_sort()`.

TBB also provides a special memory allocator, which is more efficient than the default `malloc()` implementation in multi-threaded programs.

Finally, TBB handles C++ exceptions in a special way, simplifying error management.
TBB provides a parallel sort algorithm `parallel_sort` which is based on quick sort:

```cpp
template < typename iterator >
void parallel_sort ( iterator begin, iterator end );
```

Only random access iterators are supported in `parallel_sort`, e.g. for standard C arrays or `std::vector` but not for `std::list`:

```cpp
std::vector< double > x_vector( n );
std::list< double > x_list( n );
...
parallel_sort( x_vector.begin(), x_vector.end() );
parallel_sort( x_list.begin(), x_list.end() ); // error
```

The type `item_t` of the elements to sort has to provide a swap function:

```cpp
void swap ( item_t & i1, item_t & i2 );
```

**Remark**

*The quick sort implementation of `parallel_sort` is unstable as it may change the order of identical elements, but deterministic as it will always compute the same order for the same input set.*
By default, `parallel_sort` uses `std::less` to compare elements. For user defined types or if a special ordering should be used, a third argument may be provided to `parallel_sort`:

```cpp
template < typename iterator, typename compare_t >
void parallel_sort ( iterator begin, iterator end,
                    const compare_t & compare );
```

Here, `compare_t` has to provide:

```cpp
bool compare_t::operator () ( const item_t & x, const item_t & y );
```

which should return `true` iff \( x < y \), where \(<\) is meant with respect to the corresponding ordering relation.
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```

which should return `true` iff $x < y$, where $<$ is meant with respect to the corresponding ordering relation.

The speedup of `parallel_sort` depends on the number of elements as can be seen in the right diagram. But even for large sets, the speedup is not optimal (but increasing).
Having function objects func0_t,...,func9_t implementing

```cpp
void func0_t::operator () ();
void func1_t::operator () ();
...
void func9_t::operator () ();
```

parallel execution of all functions can be accomplished by using the `parallel_invoke` algorithm:

```cpp
template < typename func0_t,
           typename func1_t,
           ...,
           typename func9_t>
void parallel_invoke ( const func0_t & f0,
                       const func1_t & f1,
                       ...
                       const func9_t & f9 );
```

The return values of all function objects is ignored and should default to `void`.

**Remark**

*The parallel_invoke algorithm is implemented for two up to ten functions.*
All functions are called by `parallel_invoke` *without* arguments. To execute functions with arguments in parallel, beside functors, lambda functions can be employed:

```cpp
void f ( int );
void g ( double );

parallel_invoke( [] { f( 1 ); },
                 [] { g( 2.1 ); } );
```

This way, also multiple return values can be implemented:

```cpp
int   f ( int );
double g ( double );
float h ( float );

int   i;
double d;
float  f;

parallel_invoke( [&i] { i = f(...); },
                 [&d] { d = g(...); },
                 [&f] { f = h(...); } );
```
Recursive execution of `parallel_invoke` is also possible:

```cpp
void traverse_tree ( node_t & node ) {
    if ( ! node.is_leaf() ) {
        parallel_invoke( [&node] { traverse_tree( node.left ); },
                        [&node] { traverse_tree( node.right ); },
        )
        handle( node );
    }
}
```
For exceptions thrown within TBB algorithms or tasks, the following handling is implemented:

1. the exception is captured and all further exception within the task group are ignored,
2. the task group is cancelled,
3. after the current algorithm has finished, the exception is rethrown by the thread which has executed the algorithm.

In order to rethrow the original exception in step 3, the program has to be compiled with C++11 support. Otherwise, the exception is of type `captured_exception`.

Assuming C++11, this enables error handling as for sequential programming:

```cpp
void main () {
    ...
    try {
        parallel_for ( ... );
    }
    catch ( my_exception & e ) {
        ...
    }
}
```
TBB comes with an extra library “tbbmalloc”, implementing optimised versions of `malloc()` and `free()`. The optimisations have a special focus for multi-threaded applications are concern

- scalability and
- false sharing.

To use it in applications, just add the library during linking:

```bash
> icpc -tbb -o myprog -c myprog.cc -ltbbmalloc
```

Furthermore, TBB implements several allocators following the C++ allocator requirements:

- **tbb_allocator**: uses the TBB versions of `malloc()` and `free()`,
- **scalable_allocator**: allocating and freeing memory scales with the number of processors,
- **cache_aligned_allocator**: allocate memory at cache line boundaries, preventing false sharing,
- **zero_allocator**: wrapper for other allocators, ensuring that the memory is zeroed after allocation.
“Intel OpenCourseWare: Threading Building Blocks”. http://intel-software-academic-program.com/courses/.

